#### An Evolutionary Framework for Matching Geospatial Object Configurations

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# Outline

- Problem Overview
- Spatial Reasoning
  - Histograms of Forces
  - Representing Object Sets
- Design of the Matching Algorithm
  - Comparing Object Sets
  - Local Search Operators
- Experiments and Results

#### **Problem Overview**

#### Where am I?

"I see a large building to my right and a small building on my left surrounded by a parking lot." Can I draw a map?



#### **Problem Overview**

#### How can I represent the sketch? How can I search the database?





#### **Problem Overview**

**Goal:** Given an input sketch and a reference database, find the most likely location from which the sketch originated.



## **Spatial Relationships**

Given a set of objects, how can we describe their spatial organization?

- Direction
- Distance
- Size
- Shape
- Topology



Should these terms be crisp or fuzzy?

A force histogram  $F_r^{AB}(\theta)$  is a way of representing the degree of truth of the statement, "*A* is in direction  $\theta$  from *B*."



(a) A force histogram is the scalar resultant of elementary forces exerted by the points of *A* on those of *B*. Each one pulls *B* in direction  $\theta$ . (b) The histogram of constant forces (r = 0) is one representation that provides a global perspective. (c) The histogram of gravitational forces (r = 2) is another possible representation, which is more sensitive to nearby points.

#### **Histograms of Forces**



#### **Comparing Force Histograms**

#### How do we compare force histograms?



Normalized Cross-Correlation:

$$\mu_{C}(f_{1}, f_{2}) = \frac{\sum_{\theta} f_{1}(\theta) f_{2}(\theta)}{\sqrt{\sum_{\theta} f_{1}^{2}(\theta)} \sqrt{\sum_{\theta} f_{2}^{2}(\theta)}}$$





#### **Normalized Cross-Correlation**

$$\mu_{C}(f_{1}, f_{2}) = \frac{\sum_{\theta} f_{1}(\theta) f_{2}(\theta)}{\sqrt{\sum_{\theta} f_{1}^{2}(\theta)} \sqrt{\sum_{\theta} f_{2}^{2}(\theta)}}$$

The normalized cross-correlation satisfies the following properties:

 $0 \leq \mu_{C}(f_{1}, f_{2}) \leq 1$   $f_{1} = f_{2} \Rightarrow \mu_{C}(f_{1}, f_{2}) = 1$   $\mu_{C}(f_{1}, f_{2}) = \mu_{C}(f_{2}, f_{1})$   $\mathbb{R}^{*} \quad \forall \lambda \in \mathbb{R}^{*} \qquad \mu_{C}(\lambda, f_{1}, \lambda, f_{2}) = \mu_{C}(f_{2}, f_{1})$ 

 $\forall \lambda_1 \in \mathbb{R}^*_+, \forall \lambda_2 \in \mathbb{R}^*_+, \qquad \mu_C(\lambda_1 f_1, \lambda_2 f_2) = \mu_C(f_1, f_2)$ 

#### **Effect of Rotation**

Rotating a pair of objects is equivalent to shifting the force histogram.



#### **Towards Rotation Invariance**

Ideally, we would like for the prior two sets to be considered equivalent.

This allows sketches to be drawn with any orientation.

This is particularly important for when hand-drawn maps are not aligned with the ground truth street grid, such as in Manhattan.



Given two pairs of objects, (A, B) and (A', B'), defined with reference angles  $\phi$  and  $\phi'$  respectively, we can compare their relative spatial relationship with the general equation

$$\begin{aligned} \mu_{Pair}(A, B, \phi, A', B', \phi') &= \beta \mu_{C0} + (1 - \beta) \mu_{C2}, & \text{where} \\ \mu_{C0} &= \mu_C \left( F_0^{AB}(\theta - \phi), F_0^{A'B'}(\theta - \phi') \right), \\ \mu_{C2} &= \mu_C \left( F_2^{AB}(\theta - \phi), F_2^{A'B'}(\theta - \phi') \right). \end{aligned}$$

 $\beta$  is a weighting factor, typically set to 0.5 to give equal weight to both histograms.

To achieve rotational invariance, we need to shift each histogram to a common value.

How far to shift?

For a force histogram  $F^{AB}(\theta)$  we define a single main direction  $\varphi^{AB}$  which best captures the directional relationship between A and B.

How to pick  $\varphi^{AB}$ ?

For a given force histogram  $F_r^{AB}$  and angle  $\theta$ , divide the forces into effective, contradictory, and compensatory forces.



Define  $b_r^{AB}(\theta)$  as the percentage of forces which are effective (maximum degree of truth).

Define  $a_r^{AB}(\theta)$  as  $b_r^{AB}(\theta)$  times the degree to which the center of mass of the effective forces aligns with  $\theta = 0$  (effective degree of truth).

[Matsakis, et al., 2001]

We should balance the contributions of the constant and gravitational forces. Define the main direction histogram as

$$\Phi^{AB}(\theta) = \max\{a_0^{AB}(\theta), \min\{a_2^{AB}(\theta), b_0^{AB}(\theta)\}\}$$



Compute the main direction  $\varphi^{AB}$  as the centroid of  $\Phi^{AB}(\theta)$ . This requires polar vector summation.



#### **Representing Object Sets**

Given a scene consisting of a set of 2D objects,  $\mathcal{O} = \{o_1, o_2, \dots, o_n\}$ , each object becomes a node in an attributed relational graph (ARG) and the relationships between objects become edges.



 $\mathcal{O} = \{o_1, o_2, o_3\} \qquad E_{\mathcal{O}} = \{e_{12}, e_{13}, e_{21}, e_{23}, e_{31}, e_{32}\}, \qquad \text{where } e_{ij} = (o_i, o_j) \in \mathcal{O} \times \mathcal{O}$ 

Assign attributes to the nodes and edges. Node attributes are object types and edge attributes are sets of force histograms.

 $\mathcal{L} = \{\text{"building", "parking lot"}\}$  $L_{\mathcal{O}}(o_1) = L_{\mathcal{O}}(o_2) = \text{"building"}$  $L_{\mathcal{O}}(o_3) = \text{"parking lot"}$ 

$$H_{\mathcal{O}} = \{h_{12}, h_{13}, h_{21}, h_{23}, h_{31}, h_{32}\},\$$
where  $h_{ij} = (F_0^{o_i o_j}, F_2^{o_i o_j}, \varphi^{o_i o_j})$ 

Full ARG:  $G_{\mathcal{O}} = (\mathcal{O}, E_{\mathcal{O}}, L_{\mathcal{O}}, H_{\mathcal{O}})$ 



### **Graph Size**

A completely defined ARG for a set of n objects will have n vertices and  $n \times (n - 1)$  edges. For large graphs, such as the reference database, only the closest objects have edge relationships.

The size of the ARG can cut in half by only considering edges for which i < j. This is possible due to the semantic inverse property of the histograms of forces.

$$F_r^{BA}(\theta) = F_r^{AB}(\theta + \pi)$$

#### **Problem Definition**

#### Given:

Reference Set: 
$$\mathcal{R} = \{x_1, x_2, \dots, x_m\}$$
  
Sketch:  $S = (o_1, o_2, \dots, o_n)$ , where  $m \gg n$ 

#### Goal:

Find a mapping function  $\Gamma: S \to \mathcal{R}$  represented as the candidate set  $\Gamma = (x_{(1)}, x_{(2)}, \dots, x_{(n)}), x_{(i)} \in \mathcal{R}$  such that  $\Gamma(o_i) = x_{(i)}$ .

The mapping function should preserve the object labels and spatial organization of the sketch.

## Comparing Object Sets (No Rotation)

If both *S* and  $\Gamma$  are defined with the same orientation, their similarity can be defined as

$$\Psi_1(S,\Gamma) = \frac{2}{n(n-1)} \sum_{i=1}^n \sum_{j=i+1}^n \mu_{Pair}(o_i, o_j, 0, x_{(i)}, x_{(j)}, 0).$$

Both reference angles are defined as 0, resulting in no shifting.

Complexity is  $O(n^2\omega)$  where *n* is the number of objects in each set and  $\omega$  is the number of force histogram bins.

#### **Optimal Rotation Angle**

When S and  $\Gamma$  have different orientations, we rotate the sketch by some optimal angle, similar to orienting a map.

The list of angular differences represents the total mismatch between the orientations of *S* and  $\Gamma$ .

$$D = \{d_{11}, d_{12}, \dots, d_{ij}, \dots, d_{(n-1)n}\}, \qquad d_{ij} = \varphi^{o_i o_j} - \varphi^{x_{(i)} x_{(j)}}$$

Pick the optimal rotation angle  $\varphi^*$  as the median angle of *D*.

$$\varphi^{\star} = \underset{d_{uv} \in D}{\arg\min[q(d_{uv})]}, \qquad q(d_{uv}) = \pi - \sum_{i=1}^{n} \sum_{j=i+1}^{n} \left| \pi - \left| d_{ij} - d_{uv} \right| \right|$$

### **Optimal Rotation Angle**





Object Pair		Main Direction (Ground Truth)	Main Direction (Sketch)	Angular Difference	Angular Distance
и	v	$\varphi^{uv}$	$\varphi^{\ u'v'}$	$d_{uv}$	$q(d_{uv})$
Α	В	3.12	4.87	4.53	-25.39
Α	С	3.30	4.81	4.78	-27.32
Α	D	3.04	4.49	4.83	-27.49
Α	E	2.12	3.58	4.83	-27.49
В	С	3.42	4.72	4.98	-26.73
В	D	3.02	4.36	4.94	-27.03
В	E	1.74	3.18	4.84	-27.48
С	D	2.77	4.24	4.81	-27.46
С	E	1.32	2.84	4.76	-27.21
D	E	0.62	1.99	4.91	-27.21

 $\min(q) = -27.49$  $\varphi^* = 4.83 \approx -83^\circ$ 

### Comparing Object Sets (With Rotation)

The similarity of *S* and  $\Gamma$  can now be defined as

$$\Psi_2(S,\Gamma) = \frac{2}{n(n-1)} \sum_{i=1}^n \sum_{j=i+1}^n \mu_{Pair}(o_i, o_j, \varphi^*, x_{(i)}, x_{(j)}, 0).$$

The sketch is rotated by the optimal rotation angle and the reference set remains unchanged.

Rather than rotate the entire sketch by the same angle, each force histogram can be shifted individually.

This allows for small imperfections between the two object sets.

Each pair of force histograms is aligned by their main direction and multiplied by a weighting factor, based on how well they correspond to the previously calculated optimal rotation angle.



$$\Psi_{3}(S,\Gamma) = \frac{2}{n(n-1)} \sum_{i=1}^{n} \sum_{j=i+1}^{n} \mu_{Trap} (\varphi^{\star} - d_{ij}) \times \mu_{Pair} (o_{i}, o_{j}, \varphi^{o_{i}o_{j}}, x_{(i)}, x_{(j)}, \varphi^{x_{(i)}x_{(j)}})$$







Elastic angles typically increases individual fitness values.

It does this for all sets of objects, raising the overall average fitness values.





### **Reference Set Used for Testing**

The reference set used in our experiments consists of 2814 buildings and parking lots from downtown Columbia, MO and the University of Missouri campus.

The graph contains edge relationships for a maximum of 50 nearest neighbors for each object.

Objects which are too far apart or have too many other objects in between them are pruned from the graph.



## **Outline of the Search Algorithm**

- 1. Set t = 0 and create an initial population of random individuals:  $P^{(0)} = (\Gamma_1, \Gamma_2, ..., \Gamma_\eta)$
- 2. For each individual in the population,  $\Gamma_P \in P^{(t)}$ 
  - a. Generate a set of possible children through mutation:  $C = mutate(\Gamma_P)$
  - b. Select the most fit child  $\Gamma_C = \arg \max_{\Gamma \in C} \psi(\Gamma)$

c. If the child is more fit than the parent, replace  $\Gamma_P$  with  $\Gamma_C$ 

- 3. Every few generations, replace the lowest scoring individuals in  $P^{(t)}$  with new random individuals
- 4. If stopping criteria is not met, increment t and go to 2
- 5. Return top scoring individual in  $P^{(t)}$

#### **Overview of the Search Algorithm**



#### **Close-Up of the Final Convergence**



#### **Mutation Operators**

Mutation operators work as local search methods, seeking to improve the quality of each solution.

The local search is applied repeatedly to each individual in the population until a suitable solution is found.

Four mutation operators were developed:



# Single-Object Replacement (SOR)

Given a parent:  $\Gamma_P = (x_{(1)}, x_{(2)}, ..., x_{(n)})$ 

- 1. Pick an object from  $\Gamma_P$  to replace
- 2. Get the set of connected neighbors which could replace the selected object and still maintain full connectivity
- 3. For each neighbor
  - a. Create a temporary child where the seed object is replaced by the neighbor and evaluate the fitness
- 4. Return the best child

We typically run this algorithm several times with different parent permutations, using all possible replacement objects.

The complexity is then  $O(pn^3\omega K)$ .

- *p* is the number of permutations
- *n* is the number of objects in the sketch
- $\omega$  is the number of bins in each force histogram
- *K* is the maximum number of nearest neighbor connections

## Single-Object Replacement (SOR)



Goal  $\Gamma = (x_7, x_{12}, x_9, x_{11}, x_{10})$ Parent  $\Gamma_P = (x_{12}, x_{10}, x_9, x_4, x_{11})$ Permute  $\Gamma' = (x_4, x_{12}, x_9, x_{11}, x_{10})$ Pick a single object  $c_i = x_4$ Get list of possible replacements  $\mathcal{X} = \{x_2, x_5, x_6, x_7, x_8, x_{13}, x_{14}\}$ Find best replacement  $x_{best} = x_7$ Replace  $c_i$  with  $x_{best}$   $\Gamma' = (x_7, x_{12}, x_9, x_{11}, x_{10})$
# Single-Object Replacement (SOR)



Goal  $\Gamma = (x_7, x_{12}, x_9, x_{11}, x_{10})$ Parent  $\Gamma_P = (x_{12}, x_{10}, x_9, x_4, x_{11})$ Permute  $\Gamma' = (x_4, x_{12}, x_9, x_{11}, x_{10})$ Pick a single object  $c_i = x_4$ Get list of possible replacements  $\mathcal{X} = \{x_2, x_5, x_6, x_7, x_8, x_{13}, x_{14}\}$ Find best replacement  $x_{best} = x_7$ Replace  $c_i$  with  $x_{best}$   $\Gamma' = (x_7, x_{12}, x_9, x_{11}, x_{10})$ 

Given a parent: 
$$\Gamma_P = (x_{(1)}, x_{(2)}, ..., x_{(n)})$$

- 1. Pick an object from  $\Gamma_P$  as the seed object and random order for gene replacement,  $\sigma$
- 2. Create an empty solution  $\Gamma'$  and assign the seed object to  $\Gamma'(\sigma_1)$
- 3. Until  $\Gamma'$  is a complete solution
  - a. Get the nearest neighbors of the current solution,  $\Gamma'$
  - b. Pick the best neighbor for the next gene location using the same method as SOR mutation
  - c. Add this neighbor to  $\Gamma'$  and move to the next gene location
- 4. Return  $\Gamma'$  as the improved child

Usually this algorithm is run with all possible seed objects and all possible initial seed locations.

The complexity is  $O(n^5 \omega K)$ .



 $\begin{array}{c|c} x_1 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \\ x_6 \\ x_6$ 



X13



X13



X13



X13



X13

Given a parent: 
$$\Gamma_P = (x_{(1)}, x_{(2)}, ..., x_{(n)})$$

- 1. Pick two objects from  $\Gamma_P$  as the seed objects and a random order for gene replacement,  $\sigma$
- 2. Create an empty solution  $\Gamma'$  and assign the seed objects to  $\Gamma'(\sigma_1)$  and  $\Gamma'(\sigma_2)$
- 3. Until  $\Gamma'$  is a complete solution
  - a. Get the nearest neighbors of the current solution,  $\Gamma'$
  - b. Pick the best neighbor for the next gene location using the same method as SOR mutation
  - c. Add this neighbor to  $\Gamma'$  and move to the next gene location
- 4. Return  $\Gamma'$  as the improved child

This algorithm is run with all possible pairs of seed objects and all possible initial seed locations.

The complexity is  $O(n^7 \omega K)$ .











Given a parent: 
$$\Gamma_P = (x_{(1)}, x_{(2)}, ..., x_{(n)})$$

- 1. Construct the graph  $G_S$  from the sketch and initialize the fitness threshold  $\psi_{\min}$
- 2. Get the set of nearest neighbors X of the parent  $\Gamma_P$
- 3. Shuffle the order of X
- 4. Construct the nearest neighbor graph  $G_{NN}$  from  $\mathcal{X}$
- 5. Run the VF2 algorithm on  $G_{NN}$  and  $G_S$  using the fitness threshold  $\psi_{\min}$
- 6. If a match is found, return it as the improved child, otherwise decrease the value of  $\psi_{min}$  and go to 3

The time complexity of the VF2 algorithm ranges from  $O(n\omega K)$  to  $O(n! \omega K)$ .



Goal  $\Gamma = (x_7, x_{12}, x_9, x_{11}, x_{10})$ Parent  $\Gamma_P = (x_4, x_5, x_3, x_6, x_2)$ Build the sketch graph  $G_S = (\mathcal{O}, E_\mathcal{O}, L_\mathcal{O}, H_\mathcal{O})$ Get list of neighbor objects  $\mathcal{X} = \{x_1, x_2, \dots, x_{14}\}$ Shuffle the order  $\mathcal{X} = \text{shuffe}(\mathcal{X})$ Build the nearest neighbor graph  $G_{NN} = (\mathcal{X}, E_\mathcal{X}, L_\mathcal{X}, H_\mathcal{X})$ Run VF2 algorithm  $\mathcal{C} = \text{match}(G_{NN}, G_S)$ Return best child  $\Gamma' = (x_7, x_{12}, x_9, x_{11}, x_{10}) \in \mathcal{C}$ 



Goal  $\Gamma = (x_7, x_{12}, x_9, x_{11}, x_{10})$ Parent  $\Gamma_P = (x_4, x_5, x_3, x_6, x_2)$ Build the sketch graph  $G_S = (\mathcal{O}, E_\mathcal{O}, L_\mathcal{O}, H_\mathcal{O})$ Get list of neighbor objects  $\mathcal{X} = \{x_1, x_2, \dots, x_{14}\}$ Shuffle the order  $\mathcal{X} = \text{shuffe}(\mathcal{X})$ Build the nearest neighbor graph  $G_{NN} = (\mathcal{X}, E_\mathcal{X}, L_\mathcal{X}, H_\mathcal{X})$ Run VF2 algorithm  $\mathcal{C} = \text{match}(G_{NN}, G_S)$ Return best child  $\Gamma' = (x_7, x_{12}, x_9, x_{11}, x_{10}) \in \mathcal{C}$ 



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# VF2 Algorithm

Given two graphs  $G_1 = (V_1, E_1)$  and  $G_2 = (V_2, E_2)$  find a mapping  $M = \{(n, m) \in V_1 \times V_2 | n \text{ is mapped onto } m\}$ .

Use a state space representation in which each state is a partial mapping M(s).



 $M = \{ (n_1, m_4), (n_2, m_1), (n_3, m_2), (n_4, m_3) \}$ 

 $M(s) = \{(n_1, m_4), (n_2, m_1)\}$  $M_1(s) = \{n_1, n_2\} \quad M_2(s) = \{m_4, m_1\}$ 



[Cordella, et al., 1999]

# VF2 Algorithm

#### **VF2 Graph Matching Algorithm**

```
Procedure: Match(G_1, G_2, s)
```

**Input:** Graphs  $G_1$  and  $G_2$ 

Intermediate state s; the initial state  $s_0$  has  $M(s_0) = \emptyset$ 

#### If M(s) covers all the nodes of $G_2$ Then

**Output:** *M*(*s*)

#### Else

```
Compute the set P(s) of the pairs candidate for inclusion in M(s)

For Each p = (n, m) in P(s)

If F(s, n, m) == TRUE Then

Create a new state s' by adding p to M(s)

Call Match(G_1, G_2, s')

End If

End For

End If
```

## **VF2** Feasibility Function

The key to the algorithm is the feasibility function that determines whether or not two nodes can be added to the partial map.

$$F(s, n, m) = F_{syn}(s, n, m) \wedge F_{sem}(s, n, m)$$

 $F_{syn}(s, n, m)$  evaluates the syntactic structure of the graphs.  $F_{sem}(s, n, m)$  evaluates the semantic compatibility.

We define the semantic feasibility function as TRUE if the fitness of the partial map is greater than the threshold  $\psi_{\min}$ .

# VF2 Syntactic Feasibility Rules

 $F_{syn}(s, n, m) = R_{\text{pred}} \wedge R_{\text{succ}} \wedge R_{\text{in}} \wedge R_{\text{out}} \wedge R_{\text{new}}$ 

Look- Ahead	Rule	Condition			
_	<i>R</i> pred	Iff for each predecessor $n'$ of $n$ in the partial mapping, the corresponding node $m'$ is a predecessor of $m_{r}$ and vice versa.			
0	R <sub>succ</sub>	Iff for each successor $n'$ of $n$ in the partial mapping, the corresponding node $m'$ is a successor of $m$ , and vice versa.			
1	R <sub>in</sub>	Iff the number of predecessors (successors) of <i>n</i> that are in $T_1^{\text{in}}(s)$ is greater than or equal to the number of predecessors (successors) of <i>m</i> that are in $T_2^{\text{in}}(s)$ .			
	R <sub>out</sub>	Iff the number of predecessors (successors) of $n$ that are in $T_1^{\text{out}}(s)$ is greater than or equal to the number of predecessors (successors) of $m$ that are in $T_2^{\text{out}}(s)$ .			
2	<i>R</i> <sub>new</sub>	Iff the number of predecessors (successors) of $n$ that are neither in $M_1(s)$ nor in $T_1(s)$ (new models) is greater than or equal to the number of predecessors (successors) of $m$ that are neither in $M_2(s)$ nor in $T_2(s)$ .			

# Experiments

We evaluate the performance of the matching algorithm using resubstitution sketches taken directly from the reference set, and simplified sketches, which have been reduced to bounding boxes and given an arbitrary orientation.

#### **Resubstitution Sketches**



#### Simplified Sketches



To compare the mutation methods, 100 resubstitution sketches and 100 simplified sketches were randomly created, each containing 5 objects.

The search algorithm runs until the original sketch location is found with a maximum search time of 1000 generations. The results are averaged over 30 trials.

Mutation	Population	Replacement	Replacement
Method	Size $(\eta)$	Frequency $(\tau)$	Percent ( $\rho$ )
SOR	50	50 Generations	50%
1-Seed	50	10 Generations	80%
2-Seed	50	10 Generations	80%
VF2	10	2 Generations	80%

Mutation Method Comparison Search Parameters

Mutation	Percent Average		Median	Average	Median		
Method	Found	Generations	Generations	Time (s)	Time (s)		
SOR	51.9%	652	931	2594	2865		
1-Seed	100%	14	5	159	44		
2-Seed	99.4%	36	12	1703	494		
VF2	100%	17	9	40	14		

#### Results for Resubstitution Skatches







Fitness: 1.0

Fitness: 0.994

Fitness: 0.994

Fitness: 0.993



Fitness: 0.993





Mutation	Percent	Average	Median	Average	Median
Method	Found	Generations	Generations	Time (s)	Time (s)
SOR	49.1%	681	1000	3127	3710
1-Seed	97.0%	60	12	539	103
2-Seed	96.5%	81	21	6404	1231
VF2	94.0%	81	11	216	19

#### **Results for Simplified Sketches**



Ground Truth





Sketch



Fitness: 0.994



#### Fitness: 0.958













Sketch



Fitness: 0.991







Fitness: 0.990

Fitness: 0.990

Fitness: 0.990

#### **Results for Simplified Sketches**



#### **Results for Simplified Sketches**



The first experiment suggested that the one-seed and the VF2 subgraph isomorphism mutation methods were the most promising.

To evaluate the impact of sketch size, 100 resubstitution sketches and 100 simplified sketches were randomly created, containing 4, 6, 8, 10, and 12 objects.

The search algorithm runs until the original sketch location is found with a maximum search time of 100 generations. The results are averaged over 10 trials.

Mutation	Number of	Percent	Average	Median	Average	Median
Method	Objects in Sketch	Found	Generations	Generations	Time (s)	Time (s)
	4	95.1%	22	11	89	40
	6	98.5%	12	3	207	40
1-Seed	8	99.6%	9	2	494	83
	10	94.8%	13	2	1258	150
	12	86.2%	20	2	3304	502
VF2	4	98.7%	16	9	26	10
	6	96.6%	19	11	81	24
	8	98.1%	19	11	189	62
	10	90.8%	26	13	1777	563
	12	76.5%	39	23	4986	2617





#### **Results for Simplified Sketches**

Mutation	Number of	Percent	Average	Median	Average	Median
Method	Objects in Sketch	Found	Generations	Generations	Time (s)	Time (s)
	4	80.6%	36	18	151	60
	6	95.6%	16	5	318	84
1-Seed	8	93.4%	18	5	771	194
	10	81.7%	28	7	2171	693
	12	87.2%	20	3	4317	874
VF2	4	80.4%	32	15	97	18
	6	94.3%	22	11	143	48
	8	86.0%	31	17	744	206
	10	69.2%	43	23	5927	1308
	12	78.8%	38	20	11009	3506

**Results for Simplified Sketches** 


## Impact of Sketch Size

**Results for Simplified Sketches** 



### **Real-World Example**

#### Hand-drafted Sketch





#### **Top Results**



Fitness: 0.913

Fitness: 0.912

Fitness: 0.911

Fitness: 0.906

Fitness: 0.905

# Conclusion

- The histograms of forces and ARG representation capture the relative spatial relationships between objects in a scene.
- This evolutionary framework allows a sketch to be searched for within a reference database.
- The one-seed and VF2 hybrid methods provide the best results.
  - VF2 works better for small sketches
  - One-seed is preferable for large sketches

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